## Beeps

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### The Problem



### **Email Beeps**

How do you maximize the time spent not checking email? (and instead presumably being productive)

- Email arrives at Poisson rate  $\lambda$ .
- Agent has beliefs at each time t whether an email has arrived.
- As soon the probability is at least  $p^*$  that an email is waiting, he checks.
- Email software is programmed to beep or not.

The probability that an email arrives by time t is

 $1 - e^{-\lambda t}$ 

When the beep is off this expression also represents the agent's beliefs at time t about whether there is an email waiting.

No Beep



### No Beep

When the beep is off, the agent will work for a length of time  $t^*$  given by

$$1 - e^{-\lambda t^*} = p^*$$

or

$$t^* = -\frac{\log(1-p^*)}{\lambda}$$

after which he will stop working and read email.

When the beep is on the agent learns immediately when an email arrives. Let  $\tilde{t}$  be the random time when an email arrives. Then the agents beliefs are a step function.



With the beep on the agent checks precisely when an email arrives. The expected time spent working is therefore

 $\mathbf{E}\tilde{t} = 1/\lambda$ 

Beep or No Beep?



## Beep or No Beep?

The beep should be turned on if and only if

$$t^* \leq \mathsf{E} ilde{t} \ -rac{log(1-p^*)}{\lambda} \leq 1/\lambda \ p^* \leq 1-1/\epsilon$$

- This is approximately 0.63
- If you are easily tempted by email, turn your beeper on!
- Not because you want to know when email arrives, but because you want to know when it *hasn't*.
- Doesn't depend on arrival rate of email.
- Note: assuming no discounting.

## Random Beeps

- When an email arrives, beep with probability z.
- A hybrid of the previous two.
- Conditional on a beep he checks.
- Conditional on no beep his posterior increases, but more slowly than no beeps.
- If there is no beep for a sufficiently long time he will check.
- Gives initial value

$$\frac{1-(1-p^*)^{-\frac{z}{z-1}}}{\lambda z}$$

- The payoff is not monotonic in z and typically an interior beep probability is optimal.
- In fact beep with probability 1 is never optimal.

## Other Ways To Beep

- Beep with a delay
- Beep with a random delay
- Beep with variable volume
- Sandom, delayed, variable volume beeps...

## Optimal Beeps: Model Overview

- The principal continuously observes a stochastic process.
- The agent knows the law of the process but doesn't directly observe its realizations.
- The agent forms beliefs about the state of the process an continuously takes actions.
- The principal sends messages to the agent about the state of the process.
- In order to influence the agent's action.
- The principal has commitment power.
- Dynamic extension of Kamenica and Gentzkow (2011)

### Examples

- Fending off audits by the board.
- Moral Hazard/performance reviews.
- Selling a good of fluctuating/depreciating quality...

### The Stochastic Process

- Finitely many states  $s \in S$ .
- Conditional on being in state s<sub>t</sub>,
  - State transitions occur at Poisson rate  $\lambda \ge 0$ ,
  - When a transition occurs the new state is drawn from a distribution

$$M_s \in \Delta S$$
.

- If  $\mu_t \in \Delta S$  is the probability over states at time t, then
  - In continuous time,

$$\dot{\mu}_t = \lambda \left( M_\mu - \mu \right)$$

where  $M_{\mu} = \sum_{s} \mu(s) M_{s}$ .

• In discrete time,  $\mu_{t+1} = f(\mu_t) := \mathbf{P}_{t+1}(\cdot \mid \mu_t)$ 





The agent has prior belief  $\mu_0$  about the initial state  $s_0$ .





 $\mu_0$  $\mu_t$ 

### In each period *t*, the agent enters with a belief $\mu_t \in \Delta S$ .





### The principal observes the current state $s_t$ .





The principal sends a message  $m_t \in M_t$ .





### The agent updates to an interim belief $\nu_t \in \Delta S$ and takes an action.





#### Time passes and the agent updates his belief to

$$\mu_{t+1} = f(\nu_t) := \mathbf{P}_{t+1}(\cdot \mid \nu_t)$$





The principal knows  $\mu_{t+1}$ .

The Principal's message  $m_t$  can be an arbitrary function of the complete prior history.

 $\sigma(h_t) \in \Delta M_t$ 

The Principal commits to a policy  $\sigma(\cdot)$ .

# Indirect Utility

The agent's action is chosen to maximize his payoff

$$a_t \in \operatorname{argmax}_{a} \mathbf{E}_{\nu_t} v(a, s)$$

And the Principal obtains (flow) indirect utility

 $u(a_t(v_t))$ 

Indeed we will take  $u(v_t)$  as the primitive assuming only that it is bounded and upper-semicontinuous.

### Principal's Objective

The principal chooses a policy to maximize the expectation of

$$(1-\delta)\sum_{t=0}^{\infty}\delta^{t}u(\nu_{t})$$

## Beeps

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- $S = \{0, 1\}$ . (no email, email waiting.)
- $\mu_t$ ,  $\nu_t$  are probabilities that email has arrived.

$$u(
u) = egin{cases} 1 & ext{if } 
u \leq p^* \ 0 & ext{otherwise} \end{cases}$$

- Policies
  - Beep on:  $M_t = \{0, 1\}, \sigma_t(h_t) = s_t$ .
  - Random beeps:  $\sigma_t(s_t) \in \Delta\{0, 1\}$ .
  - Variable volume beeps:  $M_t = [0, 1]$ .
  - etc.

# The Obfuscation Principle

#### Lemma

Any policy  $\sigma$  induces a stochastic process  $(\mu_t, \nu_t)$  satisfying

$$\bullet \mathbf{E}(\nu_t \mid \mu_t) = \mu_t,$$

**2** 
$$\mu_{t+1} = f(\nu_t).$$

Any stochastic process with initial belief  $\mu_0$  satisfying these properties can be generated by a policy  $\sigma$  which depends only on the current belief  $\mu_t$ and the current state  $s_t$ . The principal chooses the stochastic process  $(\mu_t, \nu_t)$  directly.

 $\bullet\,$  Given  $\mu_t,$  the natural state variable, choose a random variable  $\nu_t,$ 

• 
$$\mathbf{E}\nu_t = \mu_t$$
.

• Then 
$$\mu_{t+1} = f(\nu_t)$$
.

The principal is just telling the agent what his beliefs should be and the principal's payoff is the expected discounted value of  $u(v_t)$ .

### Discrete Time

Normalize the length of a "period" to 1 and let  $\delta = e^{-r}$ . Here is the value function

$$V(\mu_t) = \max_{\substack{\rho \in \Delta(\Delta S) \\ \mathbf{E}\rho = \mu_t}} \mathbf{E}_p \left[ (1 - \delta) u(\nu_t) + \delta V(f(\nu_t)) \right]$$

Given the current state  $\mu_t$ , the principal chooses a lottery p over his current payoff  $u(\nu_t)$  resulting in a transition to the new state  $f(\nu_t)$  and continuation value  $V(f(\nu_t))$ .

## Geometric Version

$$V = \operatorname{cav}\left[(1-\delta)u + \delta\left(V \circ f\right)
ight]$$












$$\operatorname{cav}\left[(1-\delta)u+\delta\left(V\circ f\right)
ight]$$



 $\mathsf{cav}\left[(1-\delta)u+\delta\left(\textit{V}\circ\textit{f}\right)\right]$ 



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ight]$$





$$\operatorname{cav}\left[(1-\delta)u+\delta\left(V\circ f\right)
ight]$$





#### **Fixed Point**

The value function is the unique fixed point of a contraction and therefore can be found by iteration.

## Iteration



#### Iteration



## **Optimal Policy**

Once we have a fixed point, the optimal policy can be easily inferred:

- Where V and cav  $[(1-\delta)u + \delta (V \circ f)]$  coincide:
  - Any randomization along a linear segment is optimal. (Random variable volume beeps.)
  - At a kink point the null signal is uniquely optimal.
- Where they don't coincide,
  - The value function must be linear there,
  - The optimal policy jumps to the endpoints of the linear segment.

### Continuous Time Limit



As the period length shrinks.

## Continuous Time Limit



As the period length shrinks.

### Continuous Time Limit



The continuous time limit.

Inspecting the limit value function we see that:

- Below  $p^*$  it is uniquely optimal to send no message.
- Above  $p^*$  it is uniquely optimal to randomize between  $p^*$  and 1.
- At *p*\*?

The value function is differentiable.

- Obviously below and above  $p^*$ .
- At p\*:
  - The value function is left- and right-differentiable
  - > The left-hand slope must be no steeper than the right-hand (concavity)
  - It cannot be strictly flatter.

 $V^{*}(p^{*})$ 

Differentiability allows us to calculate  $V^*(p^*)$ :

• To the left of  $p^*$  the agent is earning 1 and waiting for  $V^*(p^*)$ . This gives value

$$V^{*}(\mu) = \int_{0}^{t(\mu)} e^{-rt} dt + e^{-rt(\mu)} V^{*}(p^{*})$$

where  $\mu + (1-\mu)(1-\mathrm{e}^{-\lambda t(\mu)}) = \mathrm{p}^*$ 

- We can calculate the derivative with respect to  $\mu$ .
- To the right of  $p^*$ , the value function is linear with slope

$$\frac{-V^*(p^*)}{1-p^*}$$

- Differentiability says these must be equal at  $p^*$ .
- Yields  $V^*(p^*) = r/(r+\lambda)$ .

#### Implementation

- r/(r + λ) is the expected discounted value of a flow payoff of 1 that terminates with Poisson rate λ.
- Therefore it is also the *initial value* (i.e. starting at  $\mu = 0$ ) of the Beep On policy. A hint!
- Consider a beep with a delay of length  $t^*$  given by

$$1-e^{-\lambda t^*}=p^*$$

(Recall that this is the time it takes to reach belief  $p^*$  when the beep is off.)

## Beep With Delay $t^*$ Is Optimal.

When the policy is to beep with delay  $t^*$ ,

- Before date  $t^*$  there will be no beep for sure.
  - Thus the policy is equivalent to sending no message, as is optimal.
- At date  $t^*$  (and after),
  - His belief will be p\*.
  - There will be a beep if and only if an email arrived exactly  $t^*$  earlier.
  - These beeps arrive with Poisson rate  $\lambda$ .
  - After hearing a beep the belief moves to  $\mu = 1$ .
  - Consider the agent's updated belief after no beep.
    - **\*** He learns that no email has arrived more than  $t^*$  periods ago.
    - ★ He learns nothing about whether an email has arrived at some point later than that.
    - \* It is therefore as if he has been in a Beep Off policy, starting with belief  $\mu = 0$  for the past  $t^*$  periods.
    - **\*** His belief therefore remains pegged at  $p^*$  and he does not check email.
- Thus, at belief  $p^*$  the principal earns a flow payoff of 1 until a beep which arrives with Poisson rate  $\lambda$ . This implements the optimal value function. (Beliefs between  $p^*$  and 1 are never reached.)

# Coincidentally (?)

The optimal policy in continuous time coincides with the (a) static optimum from Kamenica-Gentzkow.

cav u

## In Fact

#### Proposition

lf u is

- Concave,
- Convex,
- Or a step function with 2 steps

then the static optimal policy is also optimal in continuous time.



Concave *u*.



 $u \circ f$  is still concave.



$$(1-\delta)\mathbf{u} + \delta(\mathbf{u}\circ\mathbf{f})$$



 $(1-\delta)u + \delta(u \circ f)$  still concave.



Convex hull of a concave function is the function itself.



etc.



Convex u.



 $(1-\delta)u + \delta(u \circ f)$  still convex.



The convex hull of a convex function is linear.



etc.

### Intuition?

- The expected value of the agent's posterior at time *t* is independent of the policy.
- Indeed it equals the posterior in the null policy (beep off.)
- So the choice of policy is just a gamble over paths whose "expected value" is the no-beep path.
- Wherever the function *u* is convex the principal is "risk-loving" over paths.
- Wherever it is concave he is "risk-averse"

Suggests that the continuous time optimum always coincides with the static optimum.
## Counterexample

$$u(\nu) = \begin{cases} 1 & \text{if } \nu = 0\\ 3/4 & \text{if } \nu \in (0, 1/2]\\ 0 & \text{otherwise} \end{cases}$$

## Ongoing Work

• Solving the model directly in continuous time, HJB equation:

$$rV = \operatorname{cav}\left[u + V' \cdot \frac{df}{dt}\right]$$

## Extensions

- Long run agent (dynamic incentives)
- Multiple agents (higher-order beliefs)
- Action affects the state
- Action generates information
- Agent has private information